
A New Generation of Fast Flexible Stream Ciphers

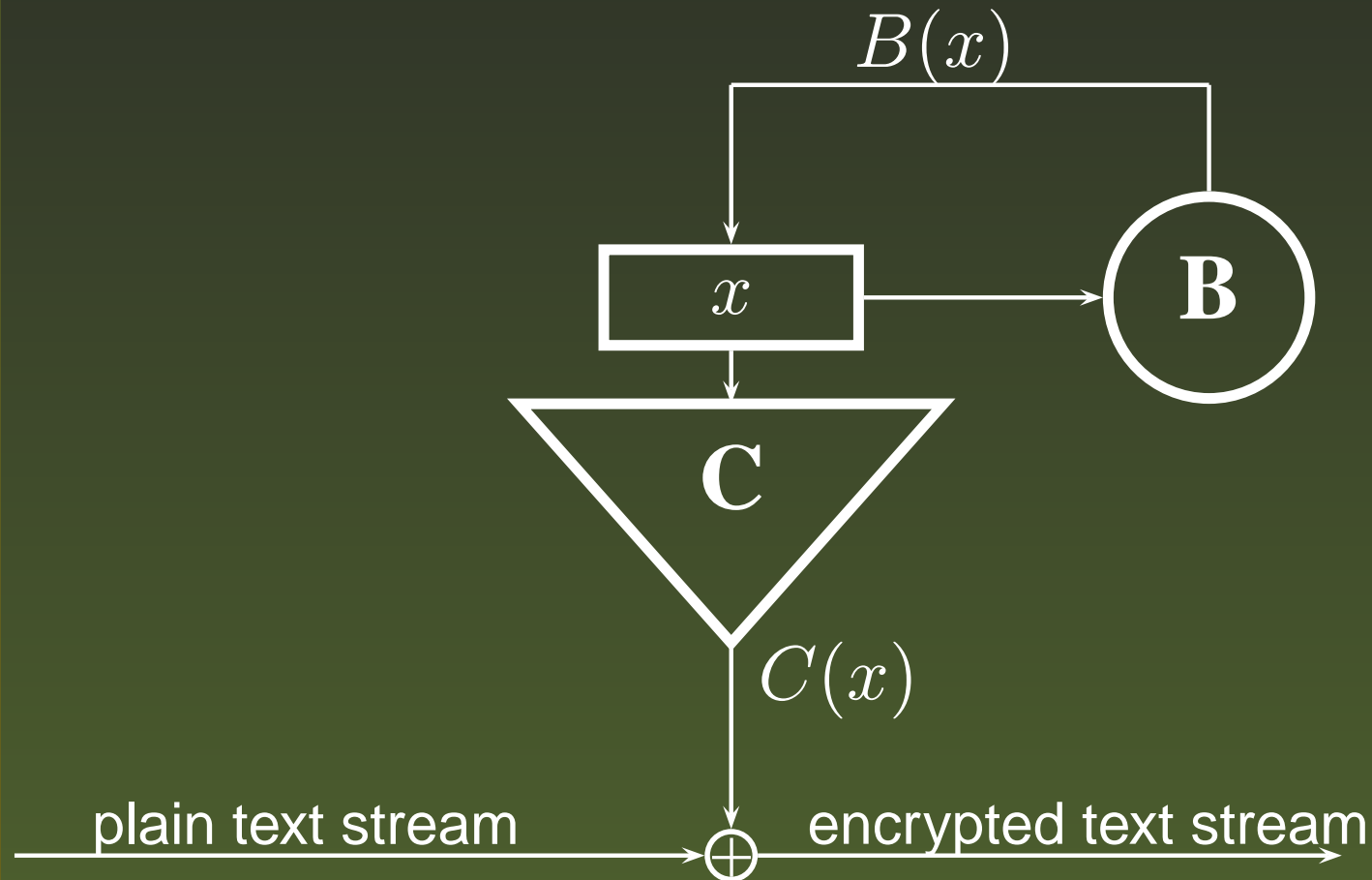
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Motivation

- Usage of new cryptographic primitives as building blocks of ciphers: T-functions and skew shifts.
- Illustration of some efficient techniques resting upon p -adic analysis and automata theory.
- A highly flexible framework for manufacturing fast and secure ciphers.

Traditional design of stream ciphers



B state transition function, period and distribution
 C non-linear filter function, other crypto properties

Traditional design of stream ciphers

Traditionally the following primitives have been used:

- State transition $B = \text{LFSR}$ over $GF(2)$ or $GF(2^k)$;
- Filter $C =$ complicated **boolean functions** or some sort of movement **irregularity**.

There are also some specific design approaches:

- R-boxes (RC4, Py, Scream, SNOW),
- S-boxes (TSC),
- Block cipher iteration procedures (SUSEMANUK, LEX, Salsa20),
- Field or ring multiplications (Rabbit),
- Feedback with carry shift registers (FCSR) and s.o.

New primitives for stream ciphers

We suggest working in rings $\mathbb{Z}/2^k$ for some suitable k and using:

- **T-functions:**
 - More simple structures to analyse;
 - First of all transitive and bijective ones;
- **Skew products** of finite automata:
 - To build complex objects from simple ones.

An efficient theory to analyse very complex encrypting mappings!

New primitives for stream ciphers

Properties:

- Long period of output keystream,
- Uniform distribution of output keystream,
- High linear complexity over $GF(2)$,
- High l -error complexity (no efficient linear analogues of keystream).

Idea:

- Make all the used maps
 - key dependent and
 - clock-dependent.

Other crypto applications

As a matter of fact the techniques are very efficient for

- block ciphers,
- MACs,
- column substitution ciphers.

All these ciphers will **inherit security properties** mentioned above from the underlying building blocks!

T-functions: Definition, Part 1

- Every map G in $\mathbb{Z}/2^n\mathbb{Z}$

$$z \in \mathbb{Z}/2^n\mathbb{Z}, z = (z_0, \dots, z_{n-1}) = \sum_{i=0}^{n-1} z_i 2^i,$$

$$G : \mathbb{Z}/2^n\mathbb{Z} \rightarrow \mathbb{Z}/2^n\mathbb{Z},$$

$$G : z \mapsto (\varphi_0(z_0, \dots, z_{n-1}), \dots, \varphi_{n-1}(z_0, \dots, z_{n-1}))$$

can be given using n (boolean) coordinate functions

$$\varphi_i : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow (\mathbb{Z}/2\mathbb{Z}), i = 0, \dots, n-1.$$

T-functions: Definition, Part 2

- $G : \mathbb{Z}/2^n\mathbb{Z} \rightarrow \mathbb{Z}/2^n\mathbb{Z}, G = (\varphi_0, \dots, \varphi_{n-1})$ is a **T-function** iff:

$$\varphi_0 = \varphi(z_0),$$

$$\varphi_1 = \varphi(z_0, z_1),$$

...

$$\varphi_{n-2} = \varphi(z_0, \dots, z_{n-2}),$$

$$\varphi_{n-1} = \varphi(z_0, \dots, z_{n-2}, z_{n-1}).$$

- T-function = triangle function.
- A number of equivalent well-known definitions in mathematics.

T-functions: Definition, Part 3

What are examples of T-functions?

- Almost **all computer operations**: $+$ and \cdot in $\mathbb{Z}/2^k$, all bitwise logical operations (AND, OR, XOR, NOT, ...), zero fill left shifts. Exceptions: Cyclic shifts, right shifts.
- All the T-functional **compositions** of T-functions **are T-functions**: Polynomials and s.o.

T-functions: Invertibility

- The T-function $G = (\varphi_0, \dots, \varphi_{n-1})$ in $\mathbb{Z}/2^n\mathbb{Z}$ is **invertible (bijective)** iff:
 - $\forall i = 0, \dots, n - 1$ the ANF of $\varphi_i(z_0, \dots, z_i)$ is linear with respect to z_i .
- It is often advisable to use invertible functions as output functions in stream ciphers.

T-functions: Single Cycle Property

- The T-function $G = (\varphi_0, \dots, \varphi_{n-1})$ in $\mathbb{Z}/2^n\mathbb{Z}$ is **transitive** (defines a single-cycle permutation) iff:
 - $\forall i = 0, \dots, n - 1$ the ANF of $\varphi_i(z_0, \dots, z_i)$ is linear with respect to z_i ,
 - the weight of φ_0 is 1,
 - for $i = 1, \dots, n - 1$ the weight of φ_i is odd,
- The linear complexity of the i -th coordinate sequence of a transitive function is $2^i + 1$. Single cycle functions guarantee the longest possible period and can be maps of choice for state update functions of stream ciphers.

T-functions: Examples of Transitive Functions, Part 1

- (Larin, early 80th; published 2002) A polynomial with integer coefficients is transitive $(\text{mod } 2^n)$ iff it is transitive modulo 8.
- (Anashin, 1993) The function $F(x) = a_0 + b_1 \cdot (x \oplus a_1) + b_2 \cdot (x \oplus a_2) + \dots$ is transitive $(\text{mod } 2^n)$ iff it is transitive modulo 4.
- (Anashin, 1993) For arbitrary polynomials $u(x), v(x) \in \mathbb{Z}/2^n\mathbb{Z}[x]$ the integer-valued function

$$F(x) = \frac{v(x)}{2 \cdot u(x) + 1}$$

is transitive $(\text{mod } 2^n)$ iff it is transitive modulo 8.

T-functions: Examples of Transitive Functions, Part 2

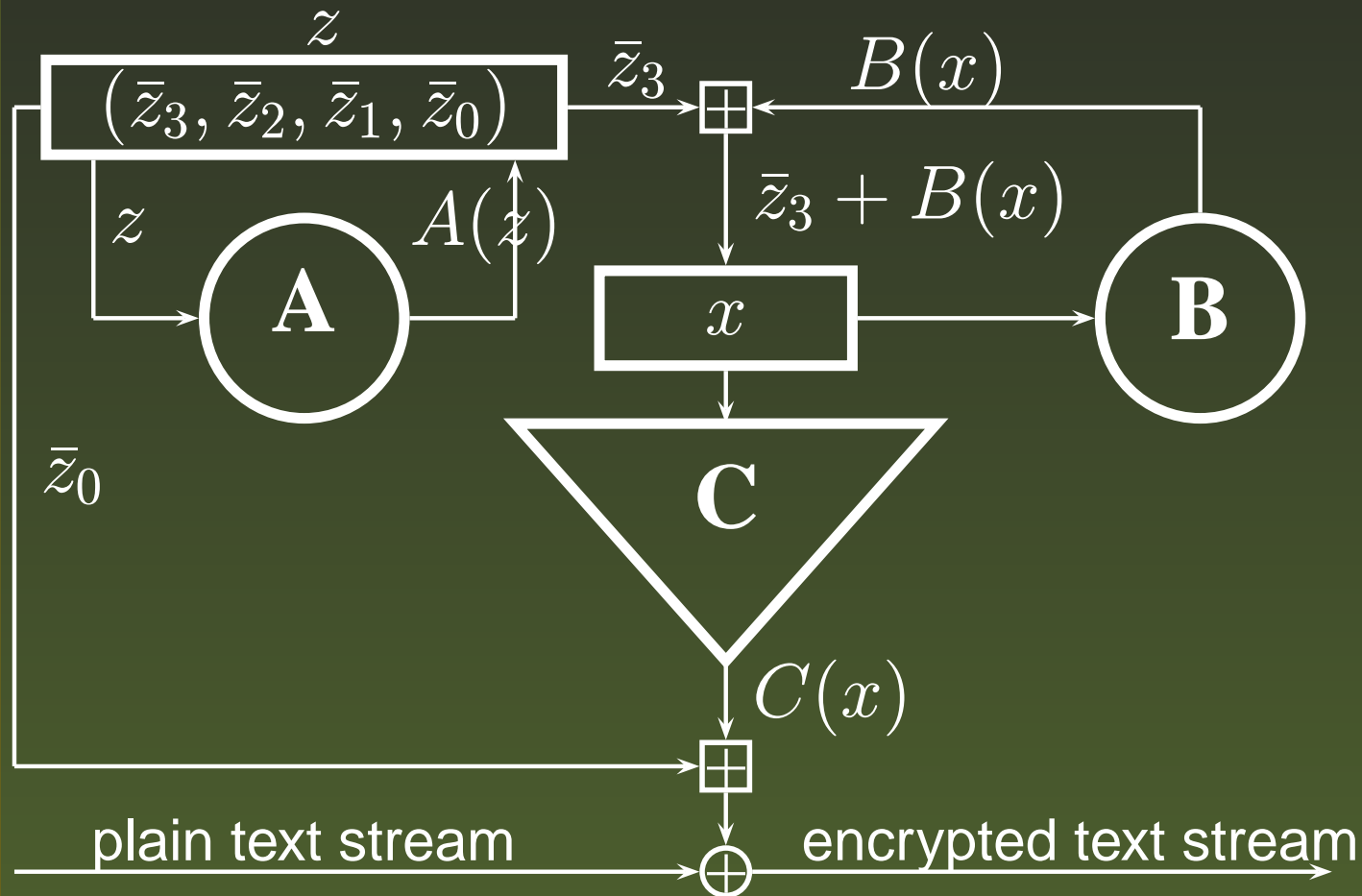
- (Kotomina, 1998) The function

$$f(x) = (\dots (((((x+c_0)\oplus d_0)+c_1)\oplus d_1)+\dots+c_m)\oplus d_m,$$

is transitive $(\text{mod } 2^n)$ iff f is transitive modulo 4

- (Anashin, 2002) The function $F(x) = a \cdot x + a^x$ is transitive $(\text{mod } 2^n)$ iff a is odd.

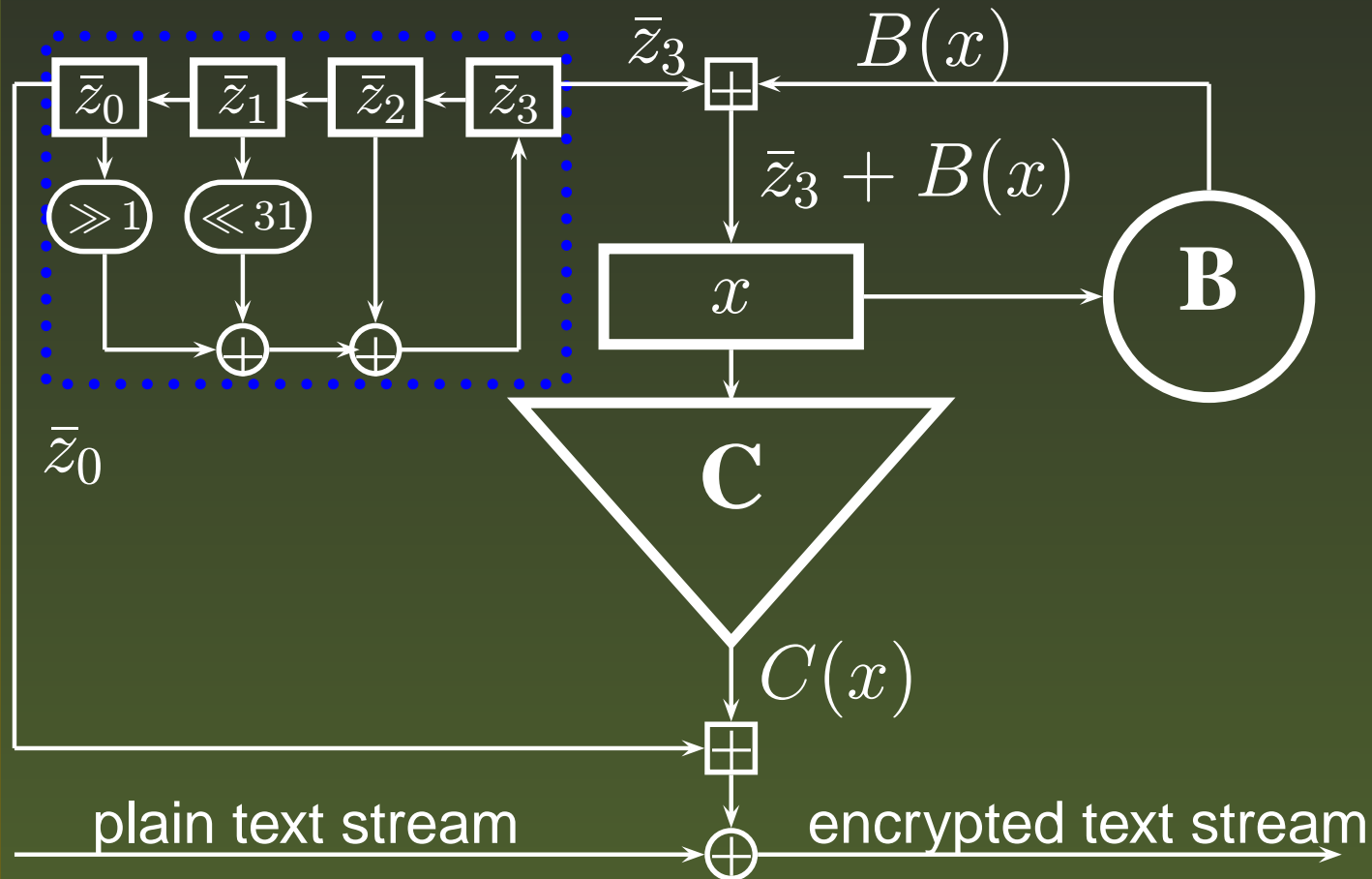
Example of new design pattern: ABC



$$\boxplus = + \pmod{2^{32}}$$

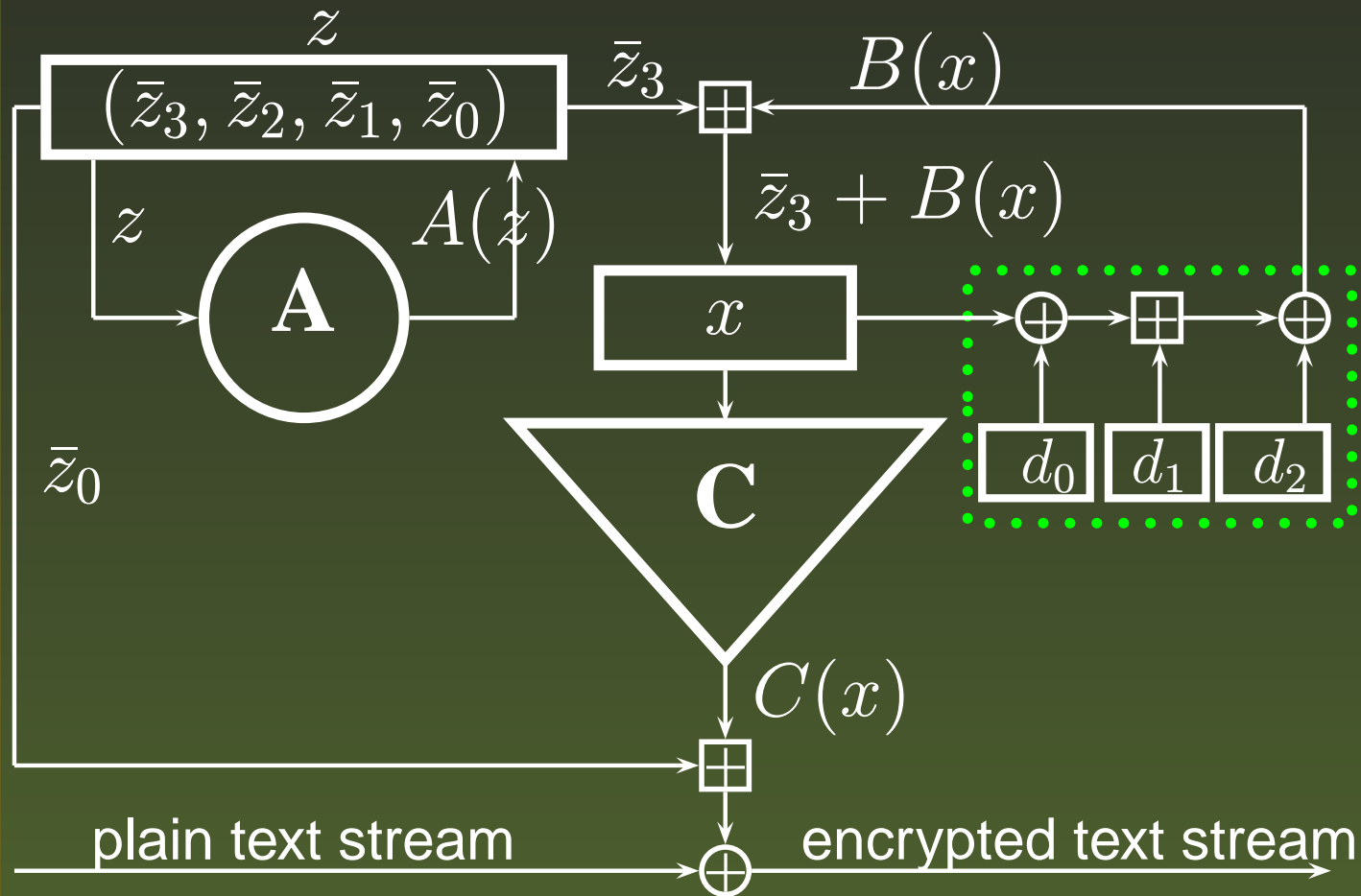
$$\oplus = \text{XOR}$$

ABC: Function A



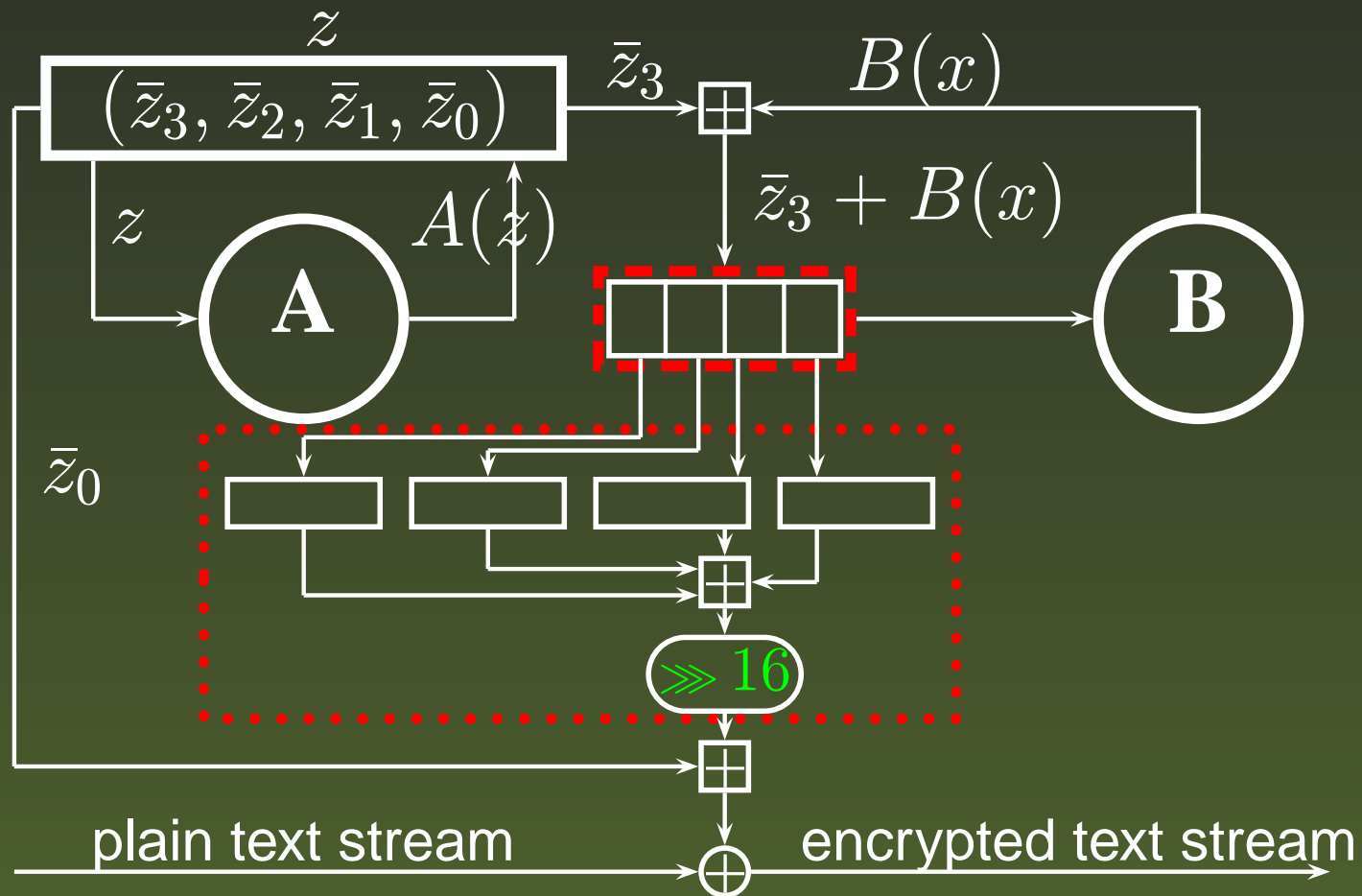
A : **LFSR** of period $2^{127} - 1$ for each 32-bit subword

ABC: Function B



B : Defines a single cycle permutation over $\mathbb{Z}/2^{32}\mathbb{Z}$

ABC: Function C



Properties of the ABC design pattern

Provable properties of the ABC key stream:

- **The period** of $(2^{127} - 1) \cdot 2^{32}$ words;
- **Uniformly distributed** key stream: \forall 32-bit word a the number $\mu(a)$ of occurrences of a on the period satisfies:

$$\left| \frac{\mu(a)}{(2^{127} - 1) \cdot 2^{32}} - \frac{1}{2^{32}} \right| < \frac{1}{\sqrt{(2^{127} - 1) \cdot 2^{32}}};$$

- **High linear complexity** λ of the key stream:
 $2^{31} \cdot (2^{127} - 1) + 1 \geq \lambda \geq 2^{31} + 1.$

Properties of ABC circuit: Notes

- As a matter of fact we have proved the group of statements for a **larger class** of A, B, C. Thus, the designer can **choose** the maps suitable for the specific requirements.
- Note that the fact that these crucial security properties are proven **does not exclude the necessity to analyse** the concrete representations of A, B and C with respect to the whole set of cryptographical attacks.

ABC: Speed & Memory consumption

- A **generic reference C** implementation on a standard 3.2 GHz Intel Pentium 4 processor under Linux.
- Minimum **132** byte memory used.

w	Speed, Gbps	Cycles per byte	Table memory, bytes
2	2.25	11.38	256
4	4.24	6.04	512
8	6.86	3.73	4096

Conclusion

- Freedom to choose mappings;
- Important security properties are *proven*;
- Novel approach to *counter-dependence*;
- High degree of *key-dependence*;
- *Key* material usage *flexibility*;
- Extremely *high throughput* rate can be achieved.